



Sistemi Embedded

A.A. 2012/2013

3CFU Automatica

Prof. Maria D. Di Benedetto

Università dell'Aquila – Centro di Eccellenza DEWS



- *Prima parte* (6 CFU ING-INF/05)
Sviluppo di hw/sw per sistemi embedded (dedicati)
- *Seconda parte* (3 CFU ING-INF/04)
Controllo embedded mediante costruzione di
modelli simbolici
(Anche parte del corso di Analisi e Controllo di Sistemi Ibridi)

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Struttura del corso:

- 24 ore
 - Costruzione di modelli simbolici per sistemi di controllo nonlineari ed ibridi
- 6 ore
 - Seminari tecnico/pratici e tesine

Modalita' d'esame:

- Prova scritta sul tema della progettazione controllori embedded “correct-by design” (10/30)

- Paulo Tabuada, Verification and Control of Hybrid Systems, A symbolic approach, Springer, 2009
- Articoli di ricerca
- Sito della didattica e-learning – Prof. M.D. Di Benedetto
- Sito del Prof. L. Pomante www.pomante.net

"Symbolic control design of
Cyber-Physical systems"

29/04/2013 - 03/05/2013

Istanbul (Turkey)

www.eeci-institute.eu



European Embedded Control Institute



Istanbul M19
29/04/2013 – 03/05/2013

Symbolic control design of Cyber-Physical systems



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Abstract of the course:

Cyber-Physical Systems (CPS) are large-scale, complex, heterogeneous, distributed and networked systems where physical processes interact with distributed computing units through communication networks. Formal approaches to the control design of these systems are relatively unexplored today. This course will present an approach to the control design of CPS based on symbolic models. Symbolic models are finite state automata where each state corresponds to an aggregate of possibly infinite continuous states and each label on the transitions to an aggregate of possibly infinite continuous inputs. We will show how the use of symbolic models provides a systematic approach to deal with control problems where software and hardware interact with the physical world through non-ideal communication networks. Efficient on-the-fly algorithms for symbolic control design will also be discussed. We will illustrate the proposed methodology on case studies.

The following topics will be covered in the course:

- * Transition systems, equivalence and compositionality
- * Approximation metrics for discrete and continuous systems
- * Incremental stability notions for nonlinear systems
- * Symbolic models for nonlinear and networked control systems
- * Symbolic control design
- * Efficient on-the-fly algorithms and case studies





Embedded Systems 2012/13



Basilica di Santa Maria di Collemaggio, 1287, L'Aquila

Lecture 1 Introduction

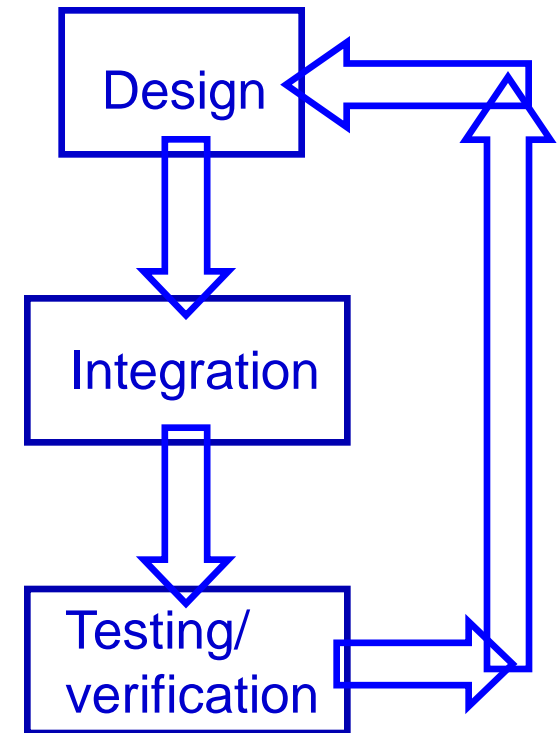
- Computational systems, but not stand alone computers
- Computation interfacing sensors and actuators
- Reactive to physical environment stimuli
- Networked and distributed information processing



Correct-by-design methodology for embedded systems design

Decomposition of concerns:

- Systems and control engineers design controllers assuming fast sampling rates for correct operation
- Software engineers design real-time operating systems and high level code assuming correct operation of controllers
- Different teams design different modules that coordinate through digital as well as physical communication
- The integrated product is then tested and/or verified to determine correctness



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Conservative design

Simulation/Verification is currently the only available method to “prove” correctness of embedded systems.
It is time and cost demanding

When a bug is found the system has to be redesigned, which may introduce new bugs

Formal verification is only possible for systems with very simple continuous dynamics, e.g. timed automata, rectangular hybrid systems, ...

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Powertrain Unit by Magneti-Marelli



Memory	256 Kb
Lines of C code	50 000
Productivity	6 Lines/Day
Changing rate	3 years
Development effort	40 man-year
Validation time	5 months
Time to market	24 months

*Fabio Romeo, Magneti-Marelli Design Automation Conference, Las Vegas, June 20th, 2001

- Electronic device controlling an internal combustion engine and a gearbox
- The goal
 - offer appropriate driving performance (e.g. torque, comfort, safety)
 - minimize fuel consumption and emissions
- Relevant characteristics
 - strictly coupled with mechanical parts
 - hard real-time constraints
 - complex algorithms for controlling fuel injection, spark ignition, throttle position, gear shift ...
 - Heuristic algorithms and simplistic average-value models used in the past

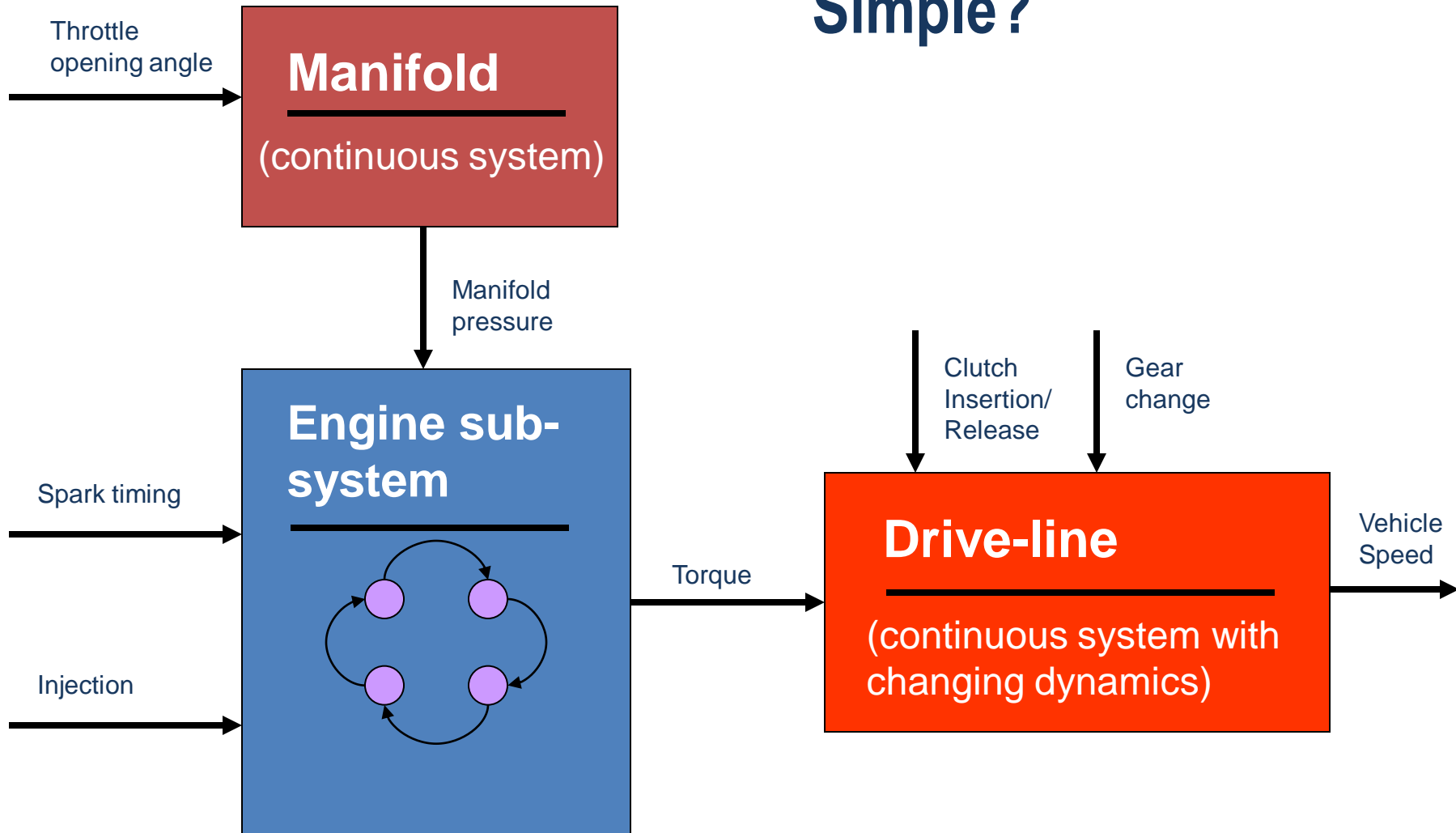
- Increasing complexity
- New safety, quality and emission requirements
- New System-on-Chip architectures
- Shorter time to market

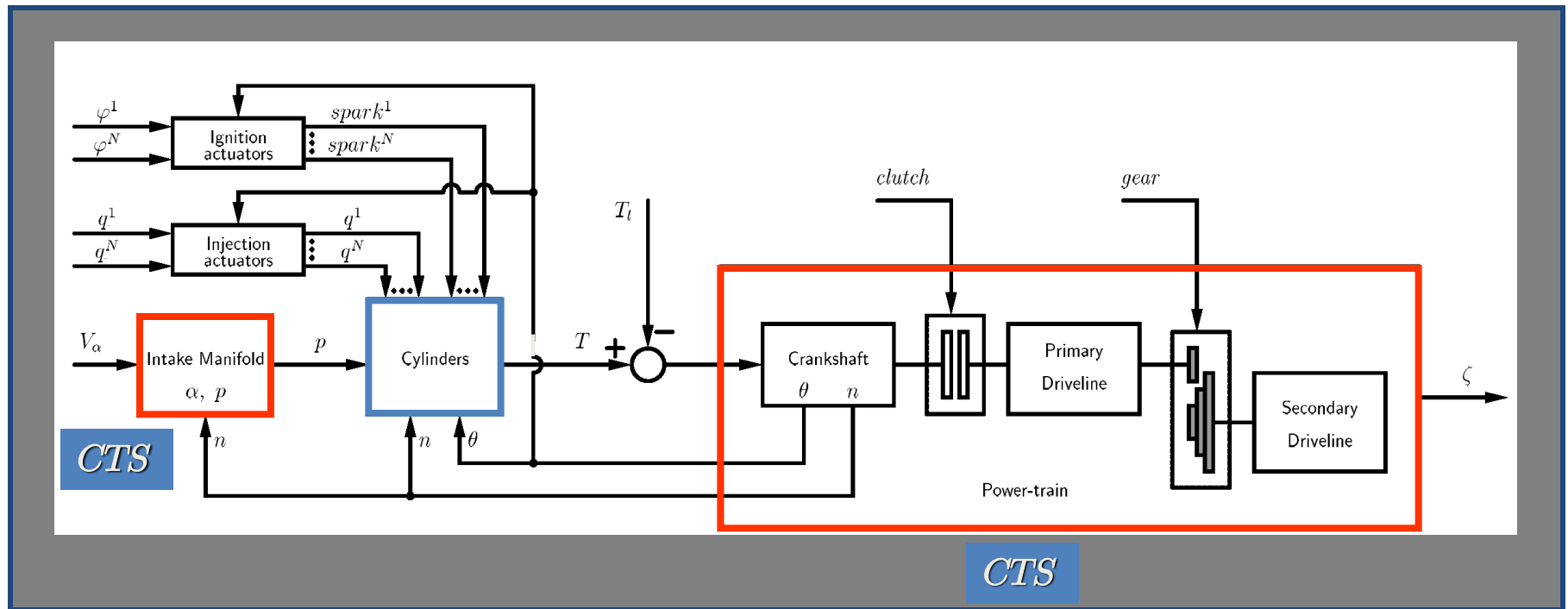


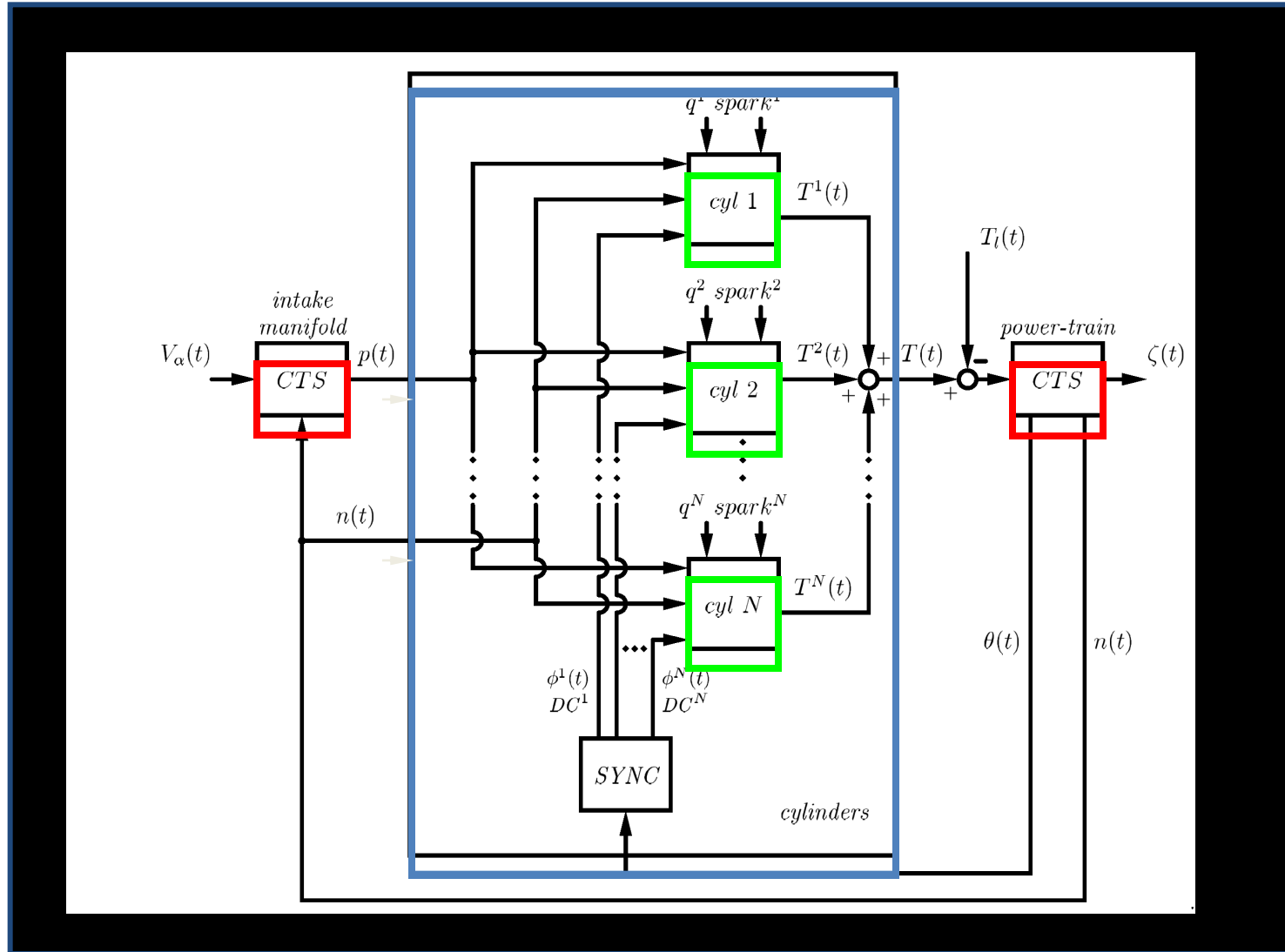
NEED more rigorous approach with

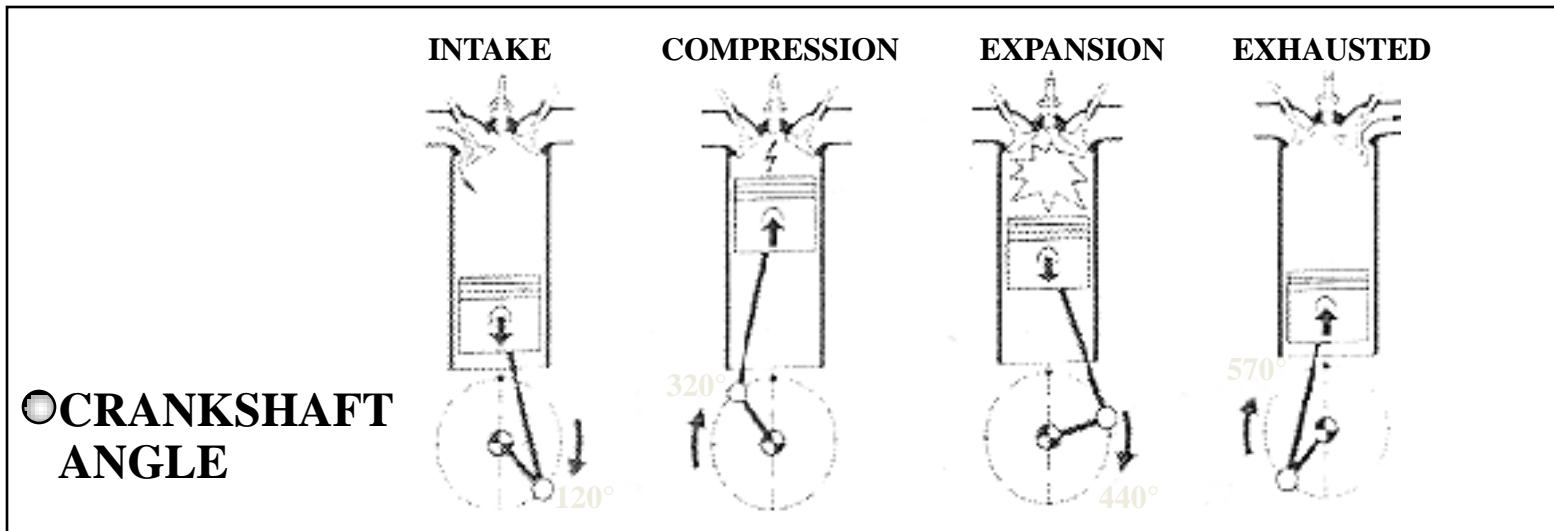
- More accurate models
- Clear partition between algorithms and implementation
- Formal proofs of correctness
- Successive refinement from specifications to implementation
- Executable Specifications

Simple?



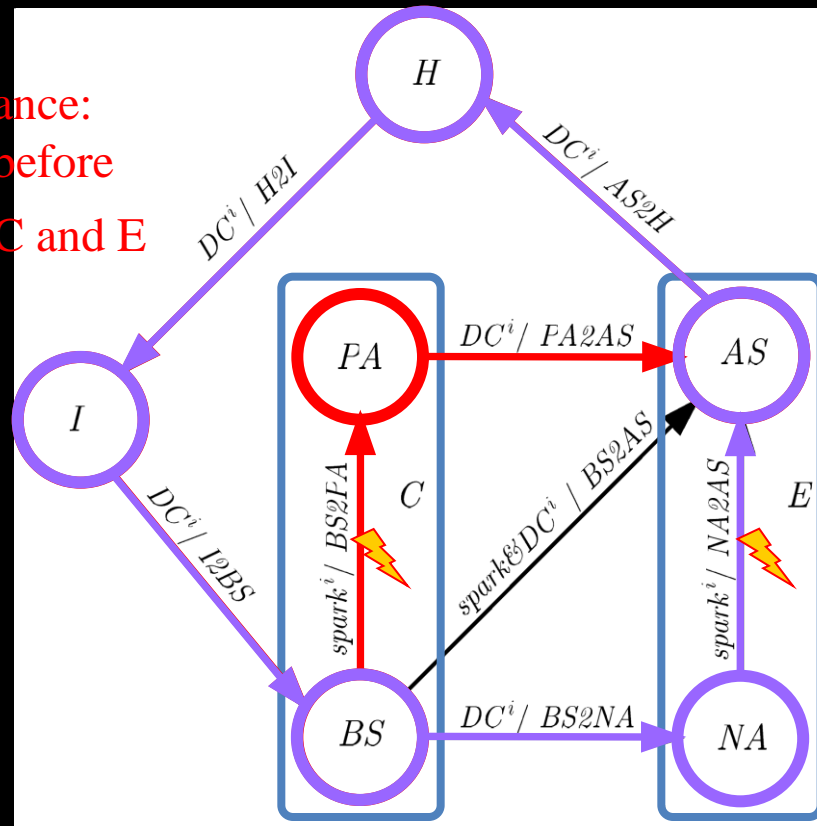


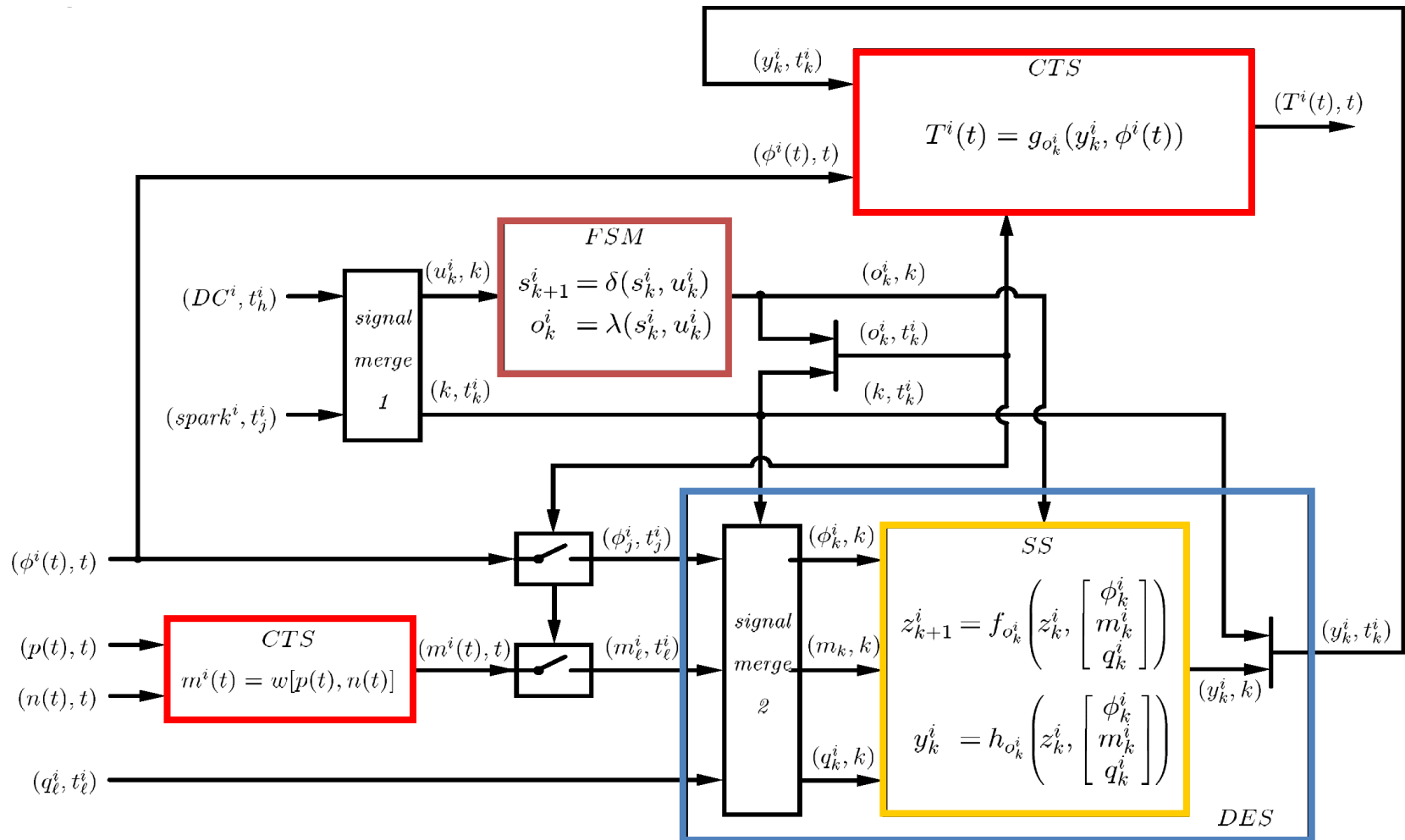




negative spark advance:
the spark is given after
the TDC between C and E

positive spark advance:
the spark is given before
the TDC between C and E





- Ariane V launched on 4th June 1996. It exploded 37s after launch
- The program had been running for 10 years, costing \$7 billions
- Software worked perfectly on Ariane IV, the same was used in Ariane V

What had changed,
was the physical
system around
the software ...



A glimpse to the future...

- 7000 billions components will serve 7 billions people in 2017 !!!

(Source: WirelessWorldResearchForum (WWRF))

- 1000 wireless components per person?





A relevant Industrial Application: Mining Ventilation Control

BOLIDEN

- Mining and smelting company
- 3 500 MEUR turnover
- 4 500 employees
- 3 Swedish and 1 Irish mine



Garpenberg mine

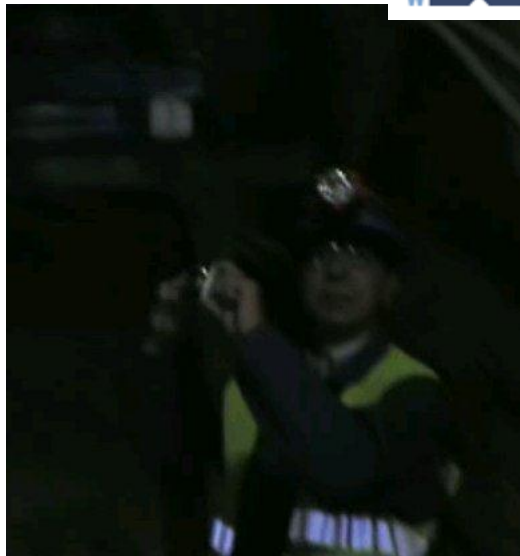
- Mining since 9th century
- 1000K ton ore/yr
 - 58K Zn, 21K Pb, 0.56K Cu, 0.1K Ag, 0.2 Au
- 1 100 m deep
- 280 employees



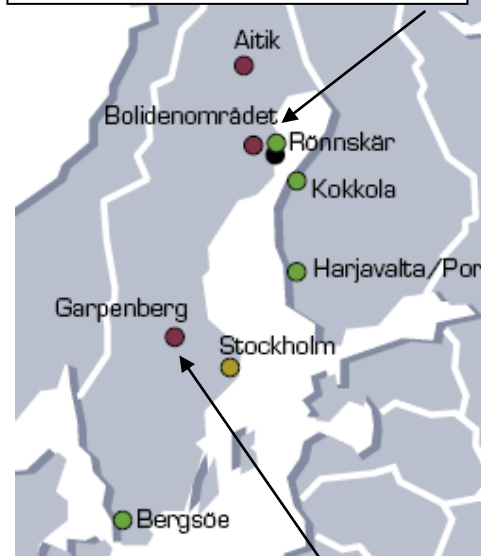
A relevant Industrial Application: Mining Ventilation Control



BOLIDEN



Boliden, 6 Mar 2007

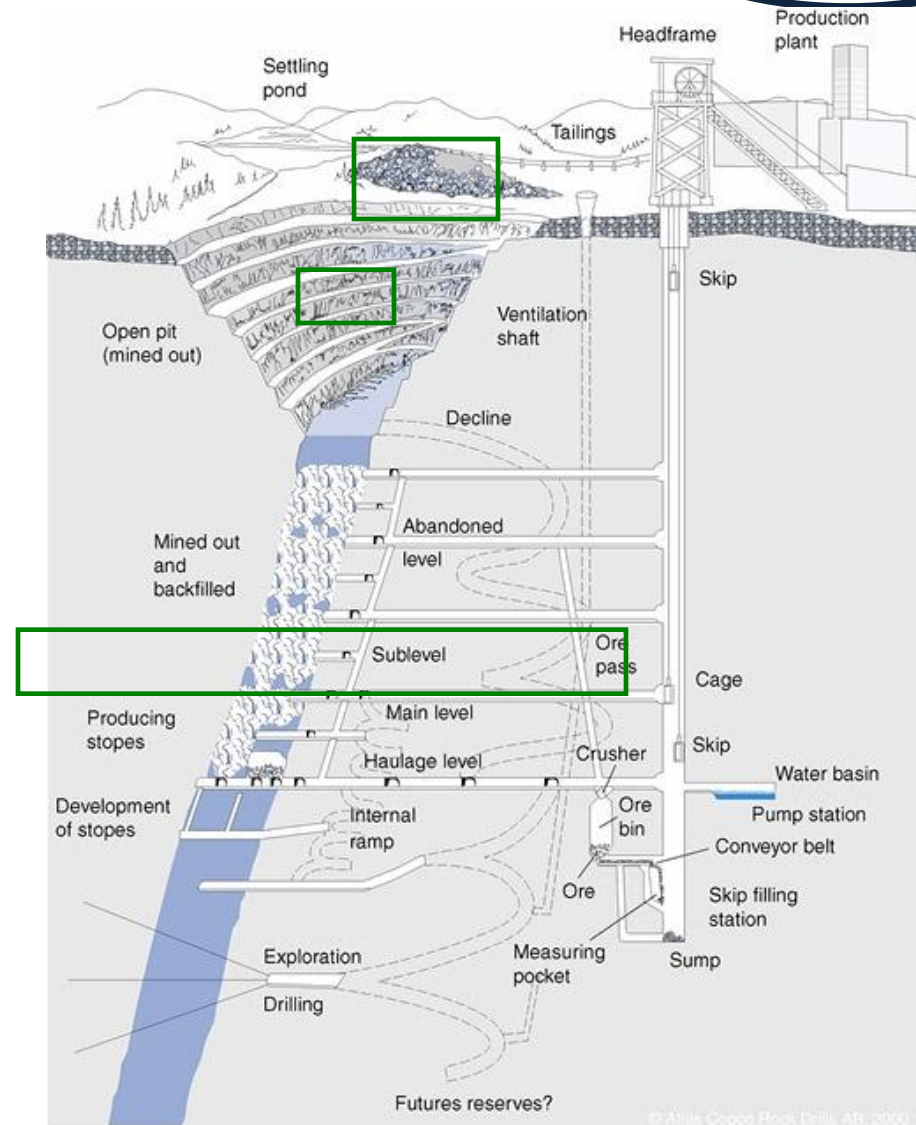


Garpenberg, 14 Sep 2007



Mining phases:

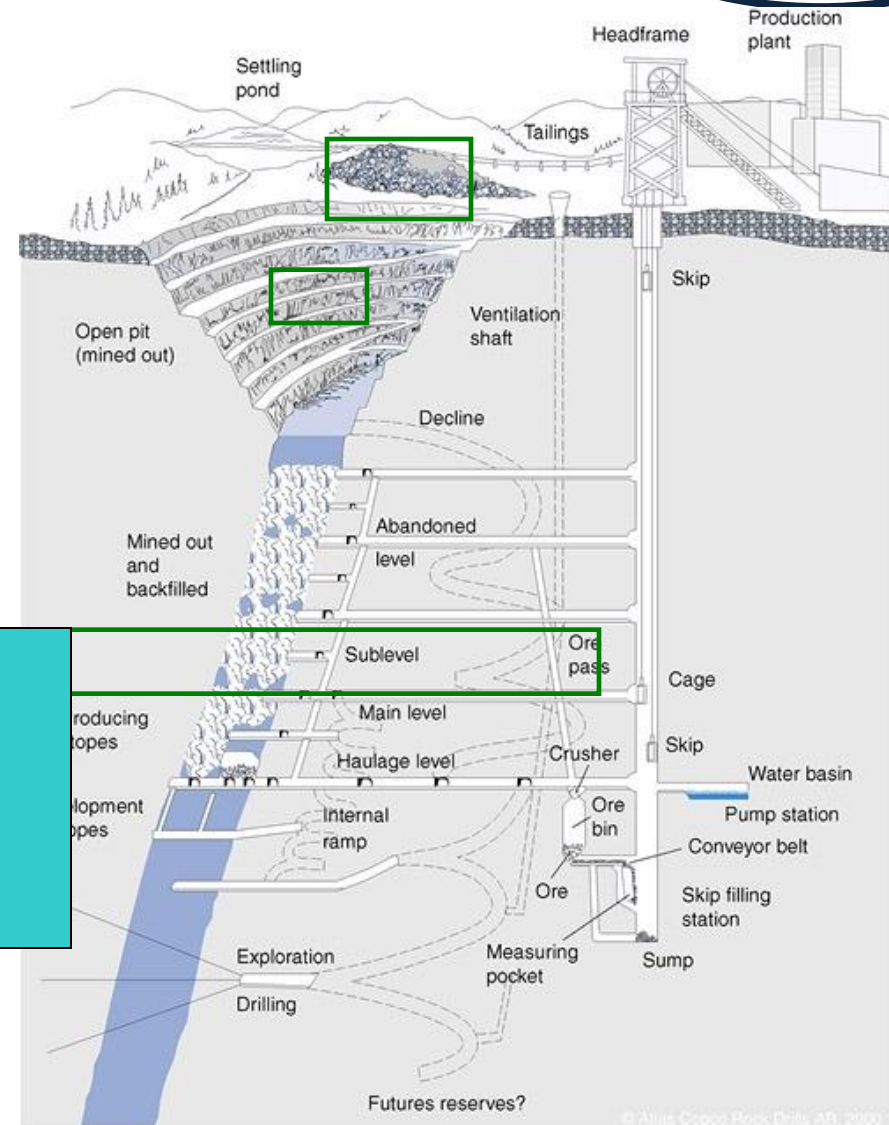
- Drilling and blasting
- Ore transportation
- Ore crushing



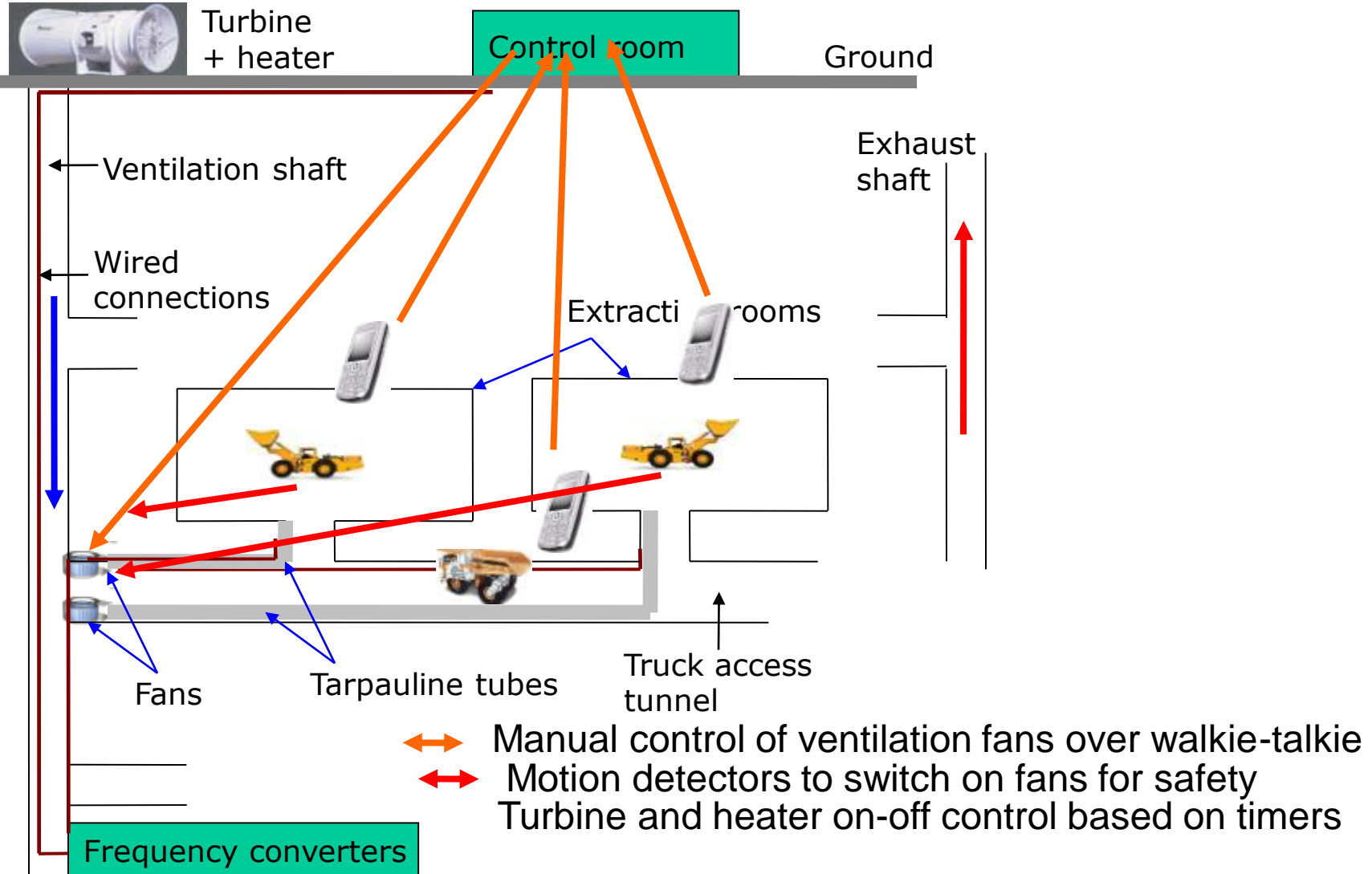
Mining phases:

- Drilling and blasting
- Ore transportation
- Ore crushing

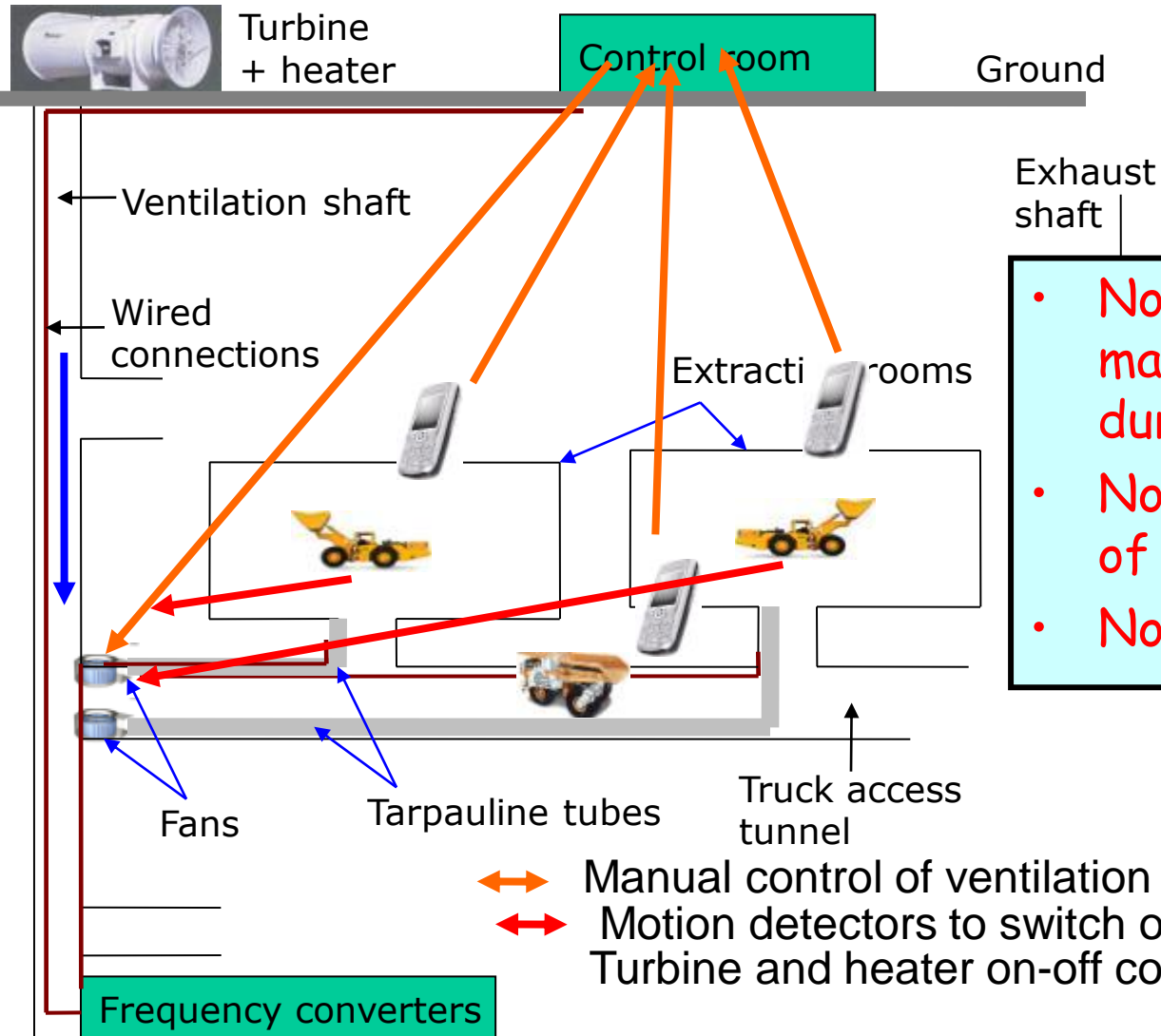
- Ventilation represents >50% of power consumption in mining
- Ventilation control system is often poor or non-existing



Mining Ventilation Control: Today's Architecture



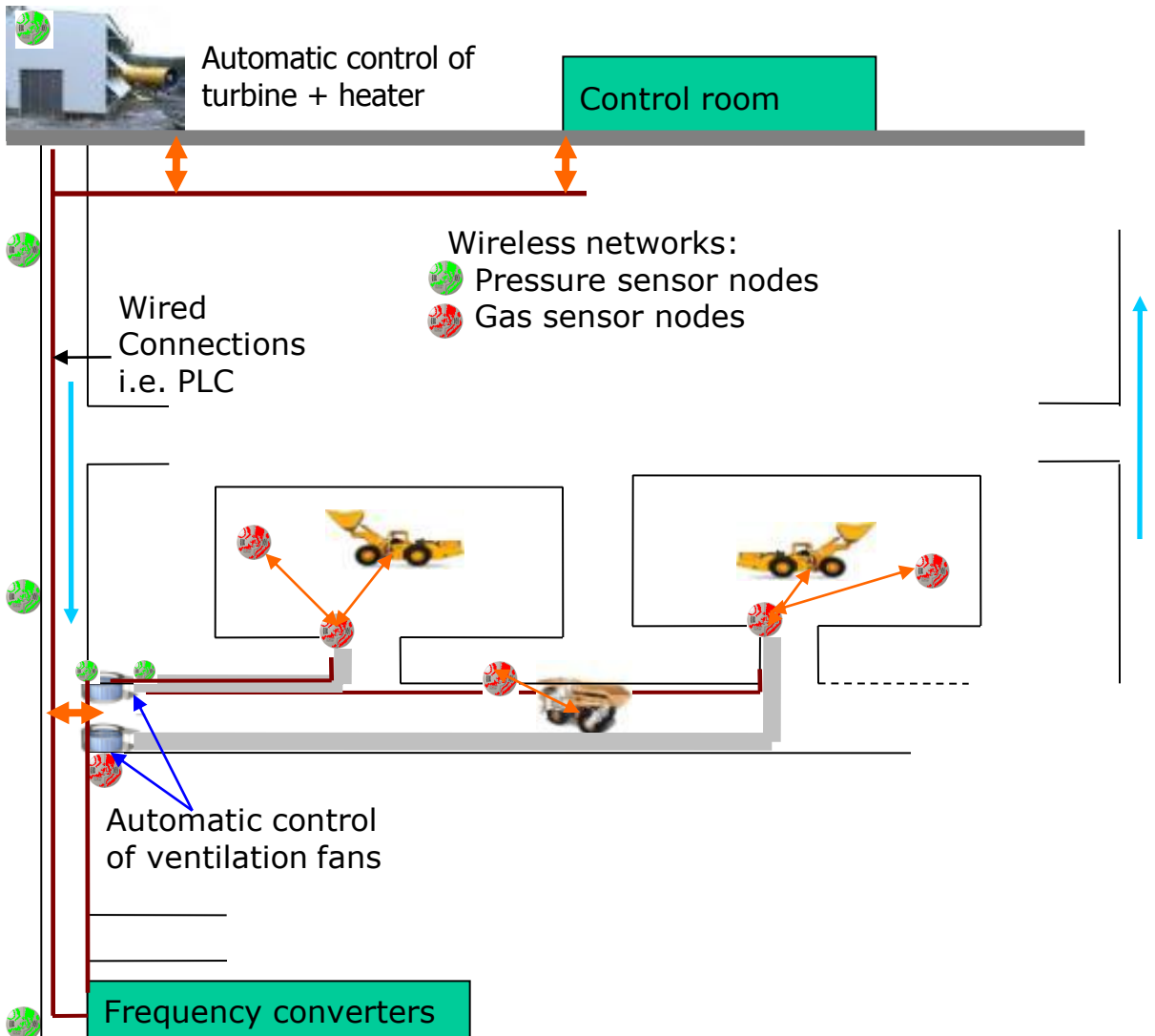
Mining Ventilation Control: Today's Architecture



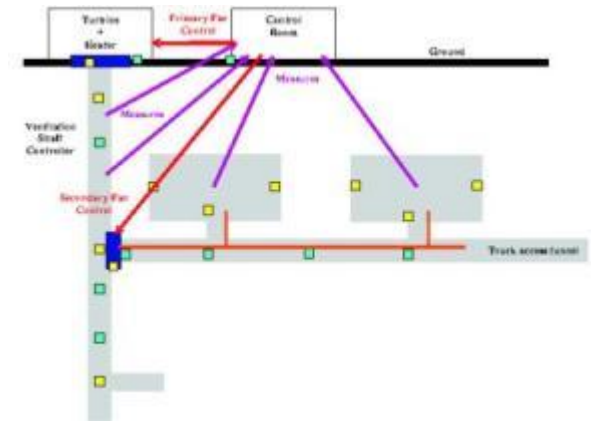
- No automatic control, but max ventilation power during ore extraction
- No continuous monitoring of air quality
- No localization system

↔ Manual control of ventilation fans over walkie-talkie
↔ Motion detectors to switch on fans for safety
Turbine and heater on-off control based on timers

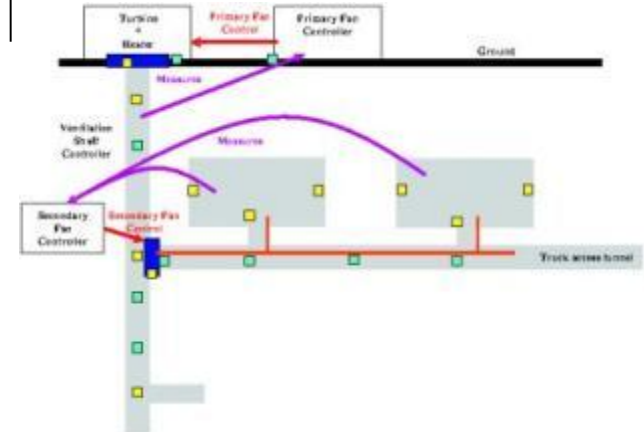
Wireless Control Architecture

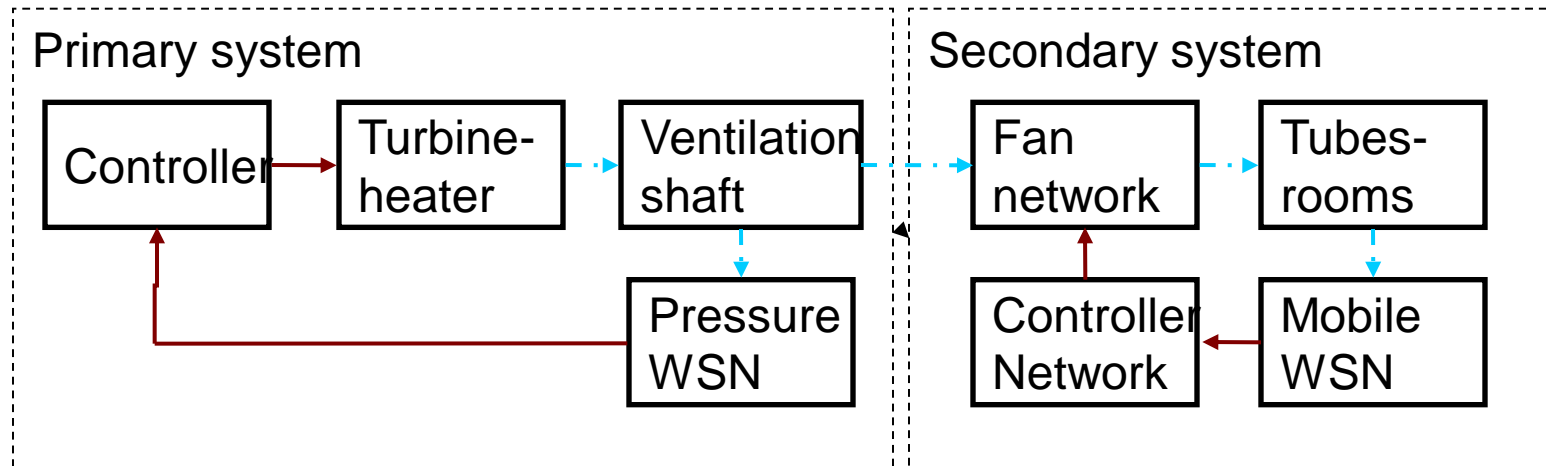


Centralized



Decentralized





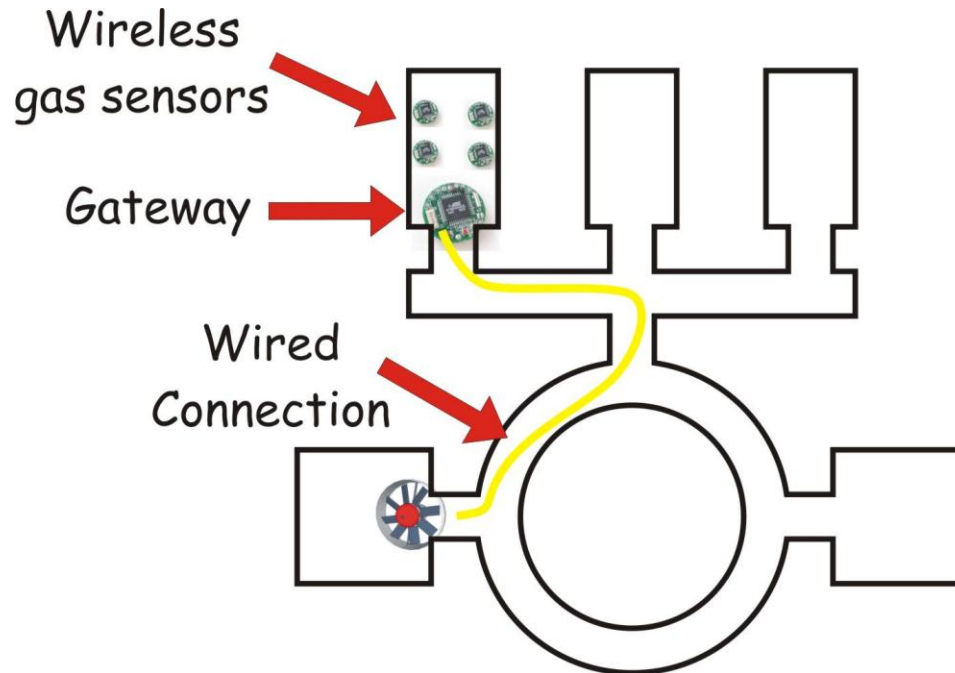
Objectives:

- Control air quality (O_2 , NO_x and CO) in extraction rooms at suitable level
 - Regulate turbine and heater to provide suitable airflow pressure at ventilation fans
 - Regulate ventilation fans to ensure air quality in extraction rooms
- Safety through wireless networking for personal communication and localization

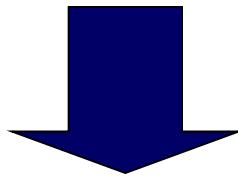
Design constraints:

- Physical interconnections, actuators limitations and networking capabilities
- Sensing capabilities: O_2 , NO_x , CO_x , pressure and temperature

- **Secondary fan:** boundary condition with primary system, rotational speed
- **Tarpaulin tube:** size, curves, junctions
- **Extraction room:** size, number of working trucks



- Secondary system may change even in the very same mine, complex geometry
- **Safety** must be **guaranteed with good margins!**



- **Our solution needs to be simple**
- **Flexible, over-approximating model**

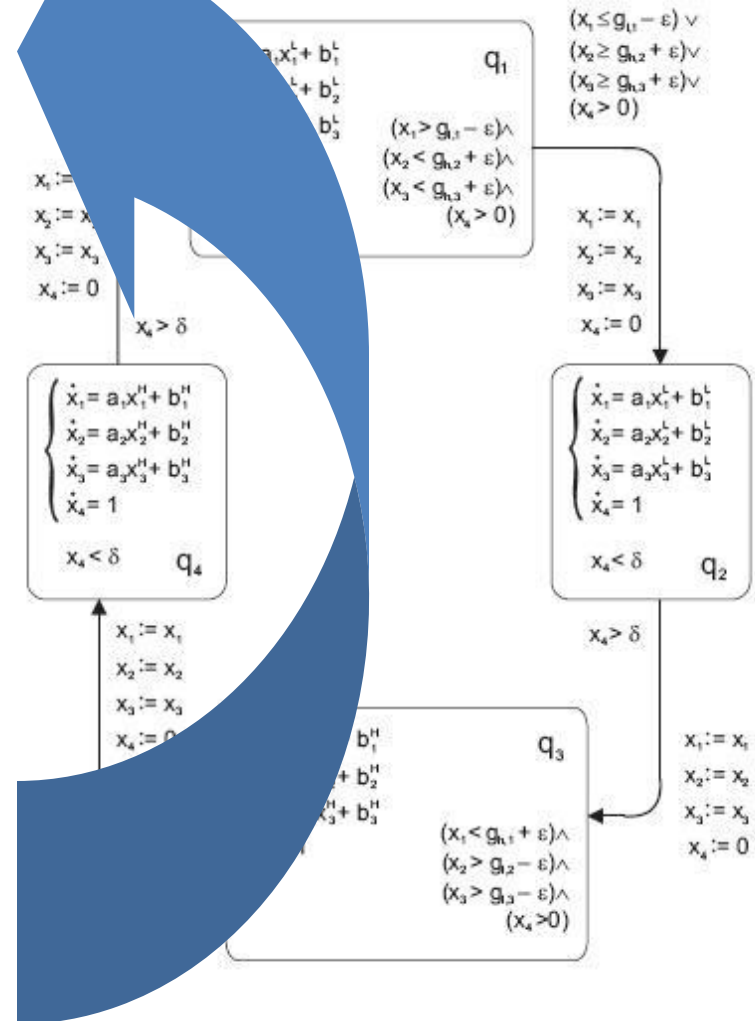
Closing the loop on wireless

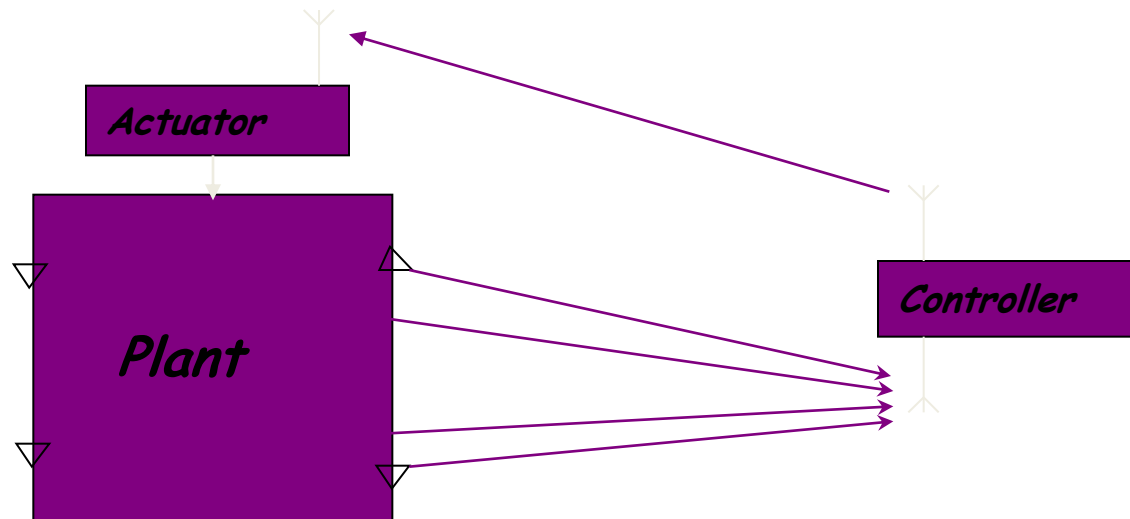
Choose simple threshold control strategy

Derive autonomous affine hybrid model of the closed loop system

Formalize safety and comfort control specifications by means of temporal logic formulae (model checking)

Perform automatic verification of the closed loop system using abstraction techniques





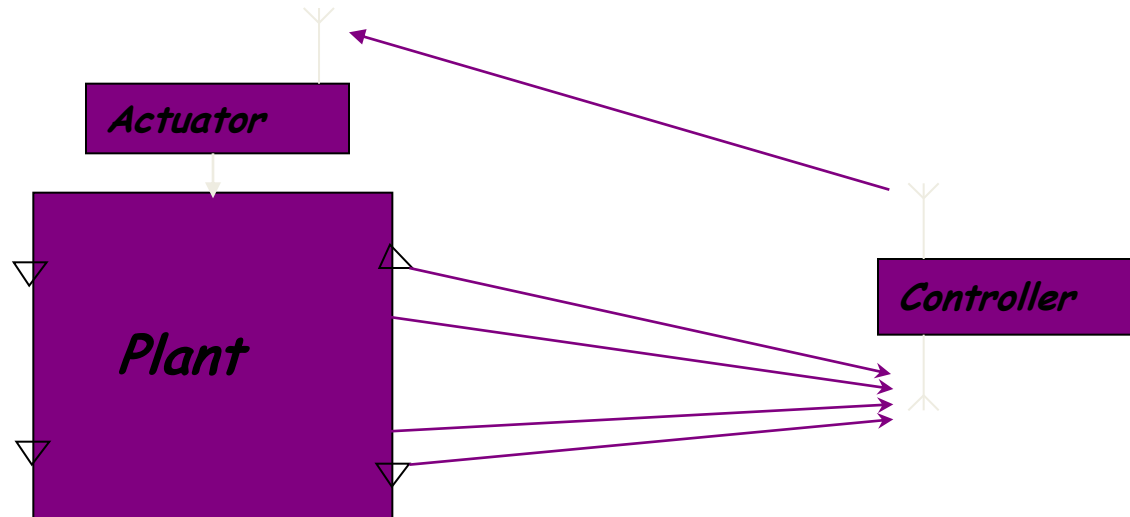
➤ Control specifications:

- Safety or Liveness
- Stability
- Robustness

➤ Wireless Network Issues:

- Packet delay
- Packet loss
- Limited spectrum
- Variable channel gain
- Interference

A communication or a control problem?



Joint Communication and Control design:

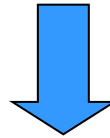
- Robustness of controller w.r.t. wireless network characteristics
- Communication scheme must satisfy constraints inferred by control specifications

Communication Systems



Control Theory

Computer Science

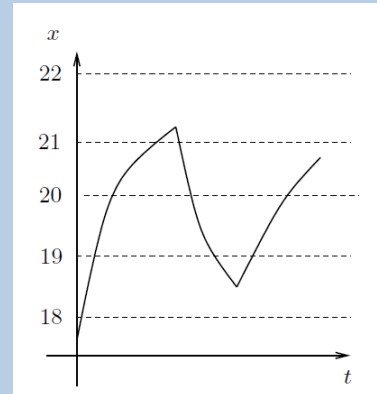
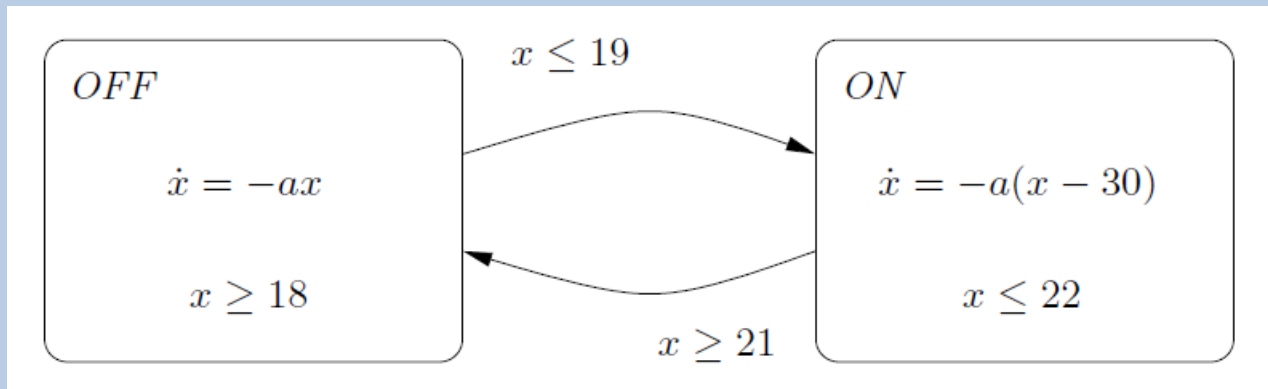


**Networked hybrid
control systems**

- Traditional approach: time-invariant systems, discrete-time and continuous-time models
- Importance of “software” and “computing systems” mixed with “physical systems”, e.g. electric circuits, mechanical control systems, etc.; these “components” form “**Cyber-Physical Systems**”
- A RLC circuit and a software are both “systems”
- In some cases, the **state** is a continuous variable in a differential equation, in other cases it corresponds to the values stored in computer memories and CPU registers
- Distributed and Networked Systems focus on communication

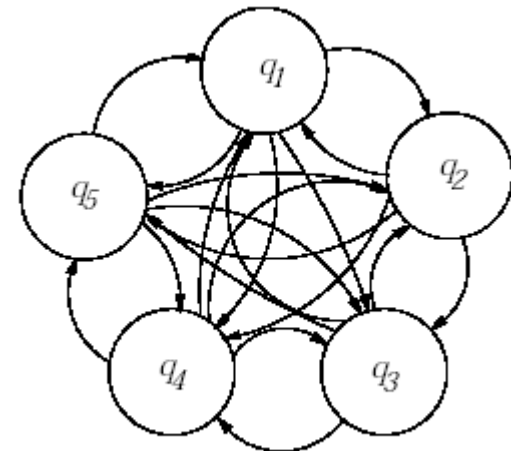
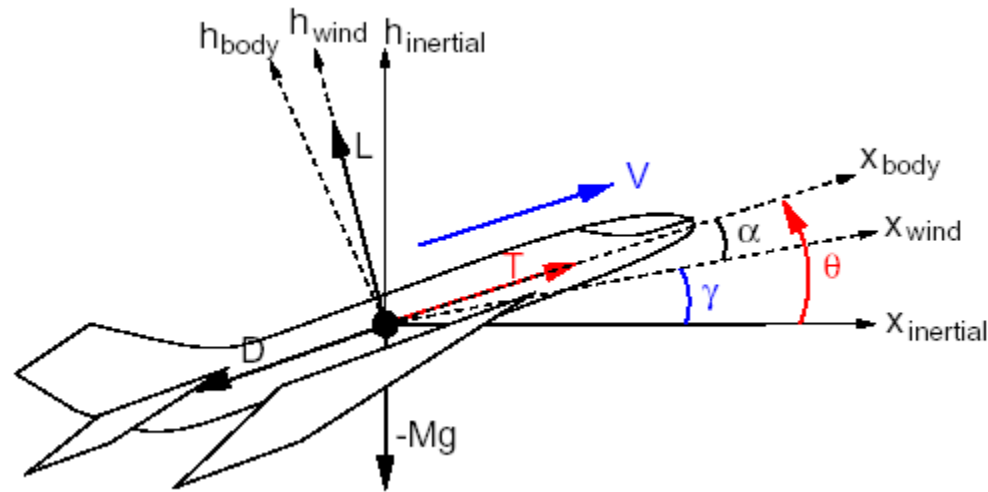
Interaction between continuous dynamical control systems and finite state machines

The Thermostat Model



From "Lecture Notes on Hybrid Systems", J. Lygeros

- **Continuous systems with a phased operation:**
 - bouncing ball
 - walking robots
 - biological cell growth and division
- **Continuous systems controlled by discrete logic:** need for design techniques that can guarantee safety and performance specifications of *embedded systems*, or systems that couple discrete logic with the analog physical environment.
 - thermostat
 - chemical plants with valves, pumps
 - control modes for complex systems, eg. intelligent cruise control in automobiles, aircraft autopilot modes
- **Coordinating processes:** many interacting subsystems (multi-agent systems) typically feature continuous controllers to optimize performance of individual agents, and coordination among agents to compete for scarce resources, resolve conflicts, etc.
 - air and ground transportation systems
 - swarms of micro-air vehicles



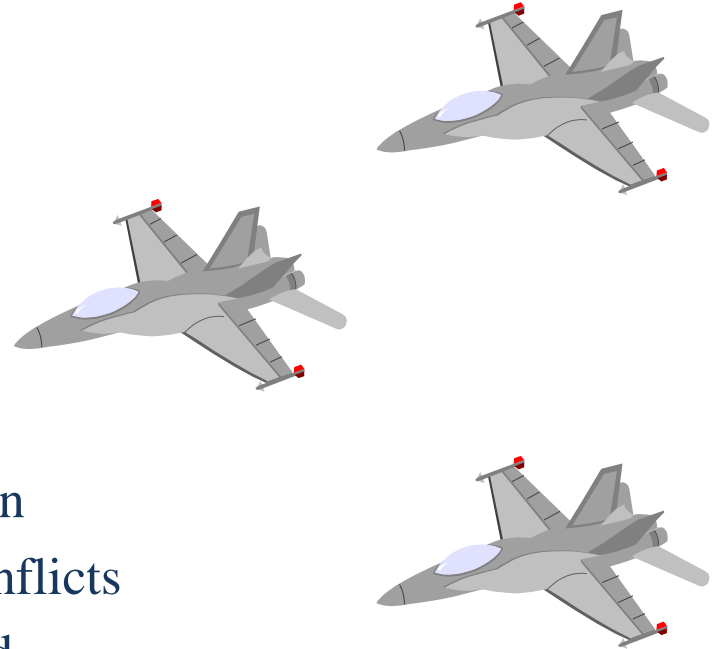
5 flight modes

Thrust T , Pitch θ

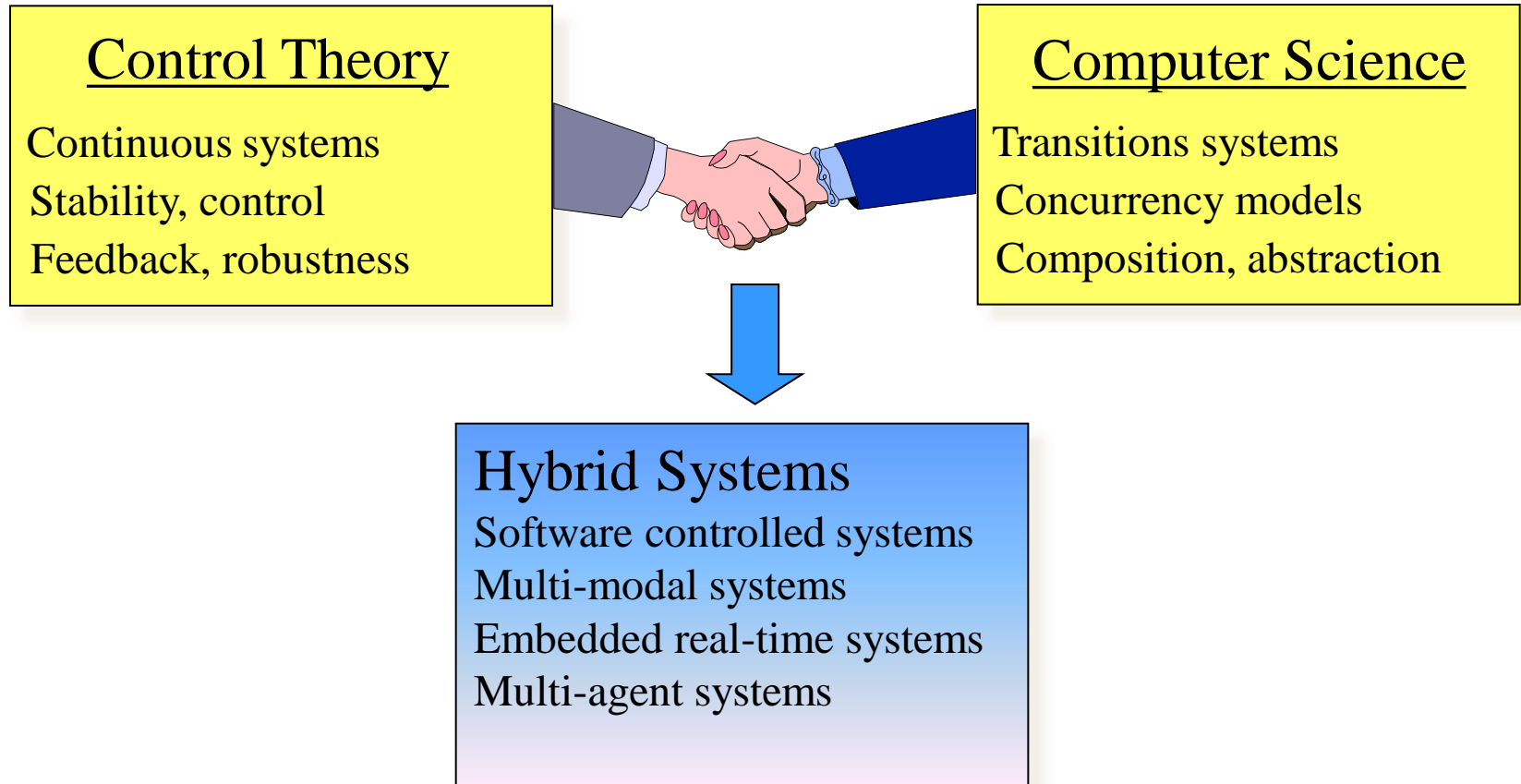
Flight path γ , Speed V

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- Large number of semiautonomous agents
- Coordinate to
 - Make efficient use of common resource
 - Achieve a common goal
- Individual agents have various modes of operation
- Agents optimize locally, coordinate to resolve conflicts
- System architecture is hierarchical and distributed
- Safety critical systems



Challenge: Develop models, analysis, and synthesis tools for designing and verifying the safety of multi-agent systems

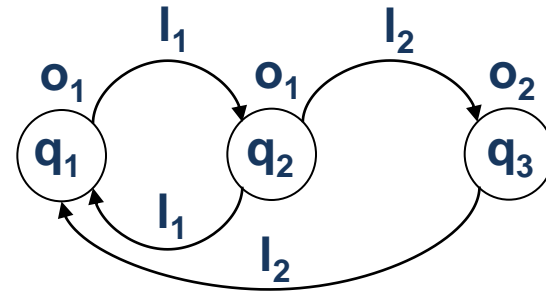


Definition A **Labelled Transition System** (for short also called **LTS** or **System**) is a tuple:

$$T = (Q, Q_0, L, \longrightarrow, O, H),$$

consisting of:

- a set of states Q
- a set of initial states Q_0
- a set of labels L
- a transition relation $\longrightarrow \subseteq Q \times L \times Q$
- a set of outputs O
- an output function $H: Q \rightarrow O$



We will follow standard practice and denote $(q, l, q') \in \longrightarrow$ by $q \xrightarrow{l} q'$

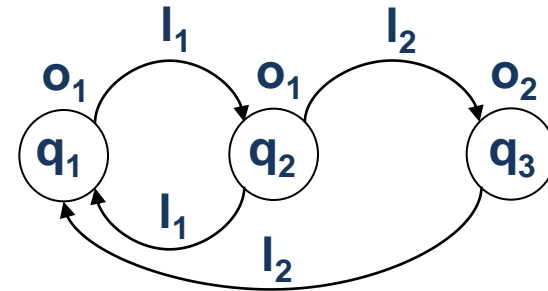
When $Q_0=Q$ we write the LTS in a more compact notation as $T = (Q, L, \longrightarrow, O, H)$

Definition A Labelled Transition System (LTS also called System) is a tuple:

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consisting of:

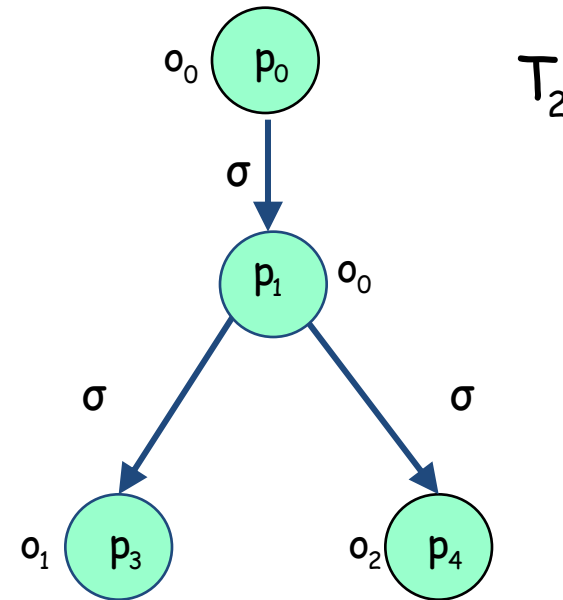
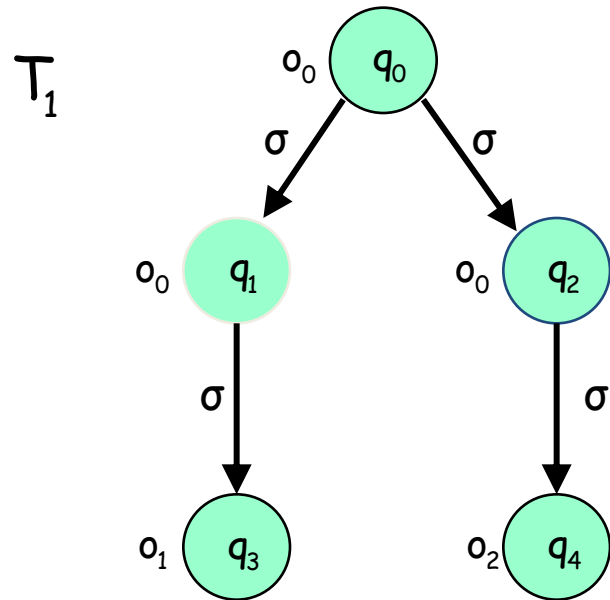
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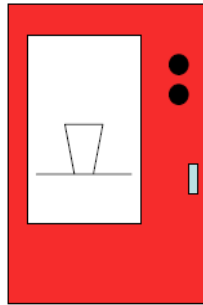
T is said:

- countable if Q and L are countable sets
- symbolic/finite if Q and L are finite sets
- deterministic if for q and l there exists at most one p such that $q \xrightarrow{l} p$
- non-blocking if for any q there exist at least one l and one p such that $q \longrightarrow p$
- metric if O is a metric space

Are T_1 and/or T_2 deterministic? Symbolic?

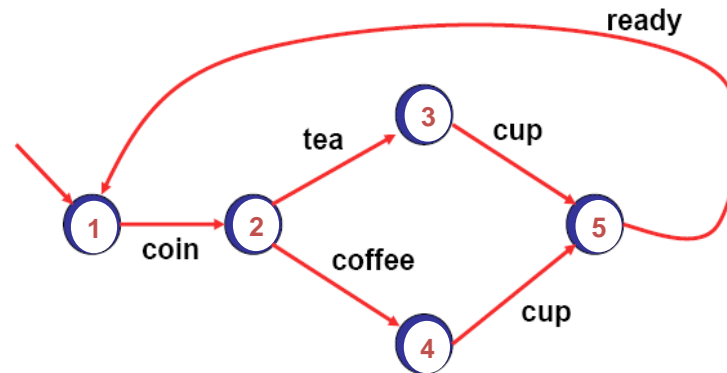


Vending machine

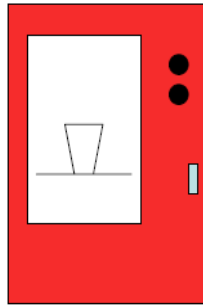


Vending machine

1. Insert coin(s)
2. Choose tea or coffee
3. Put the cup on the tray

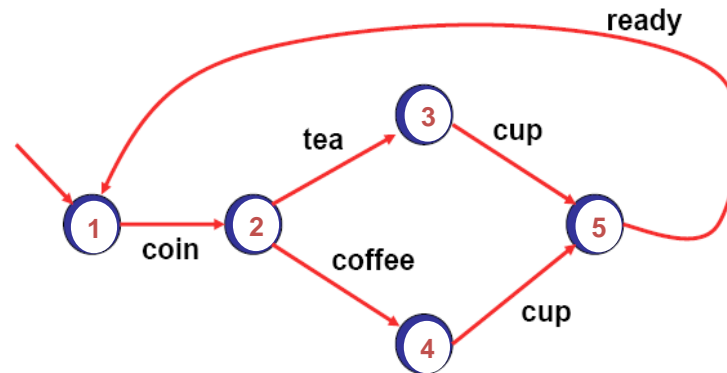


Vending machine



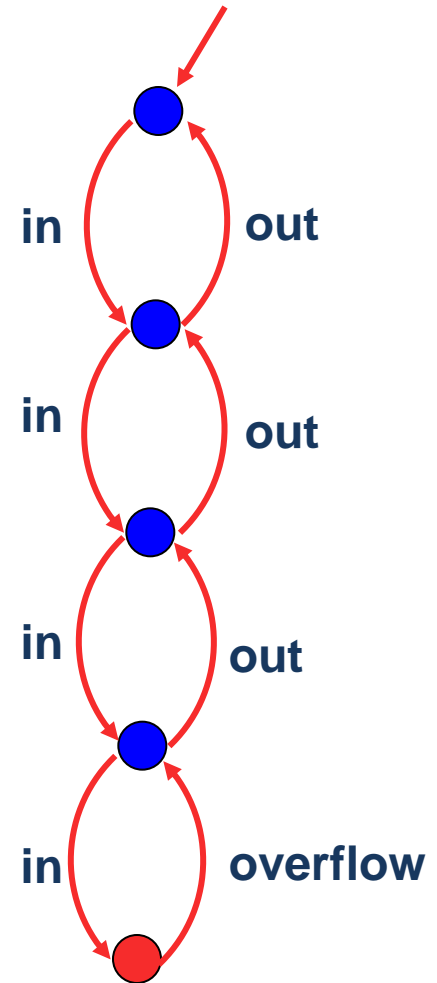
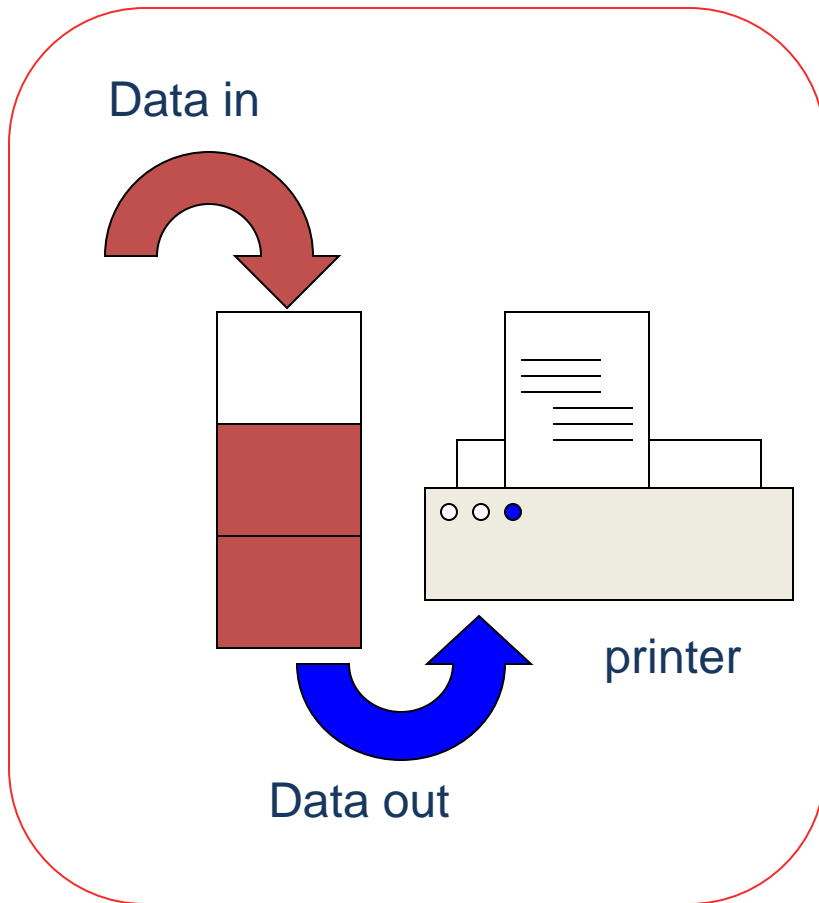
Vending machine

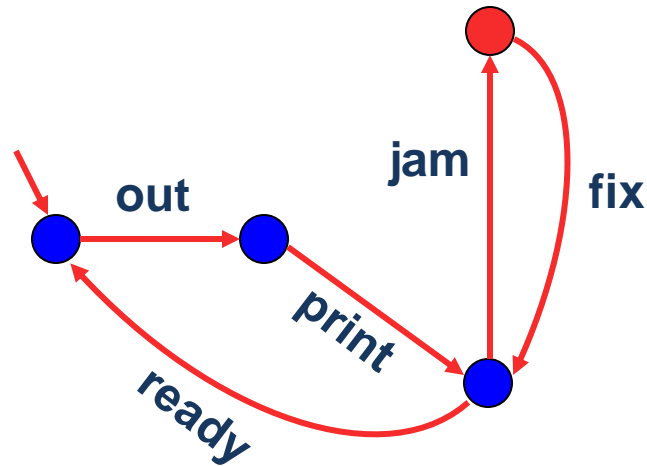
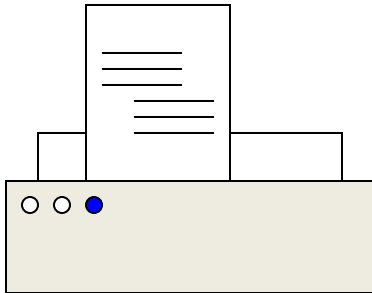
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... Event driven versus time driven !

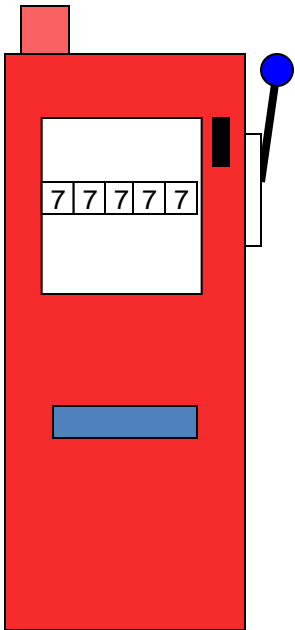
A printer data buffer



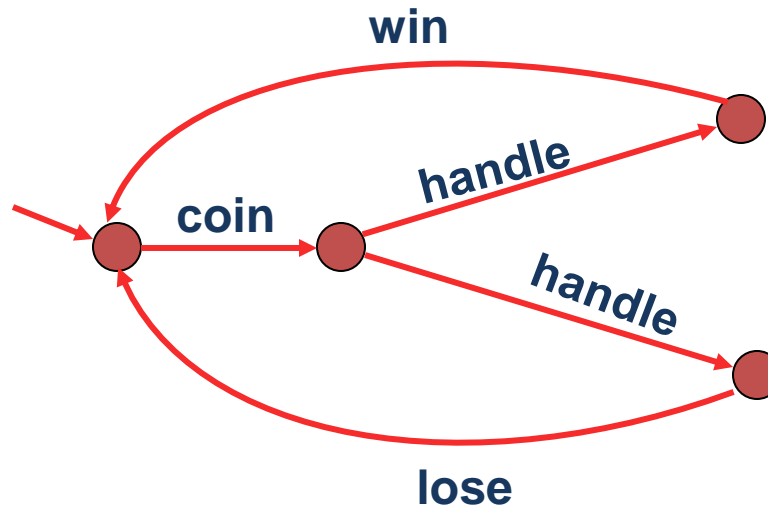


The printer receives data from the buffer, and print it out. Once the printout is ready, the printer is ready to receive new data. While printing, the paper can jam and need to be fixed before the printing process can resume.

A slot machine



1. Insert coin
2. Pull handle
3. Win if the combination is good, otherwise lose.



- Events are **time-abstract**.
- Just like modeling of continuous systems, the level of detail is '**modeler dependent**'.
- **Compositionality** is possible (to be discussed later).
- There can be **non-determinism**.

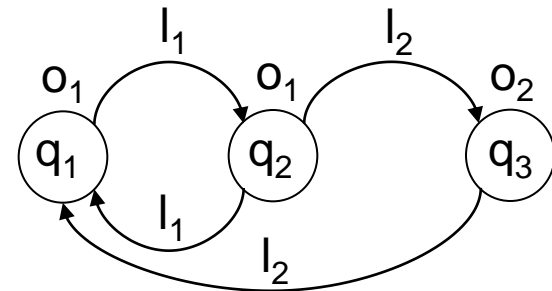


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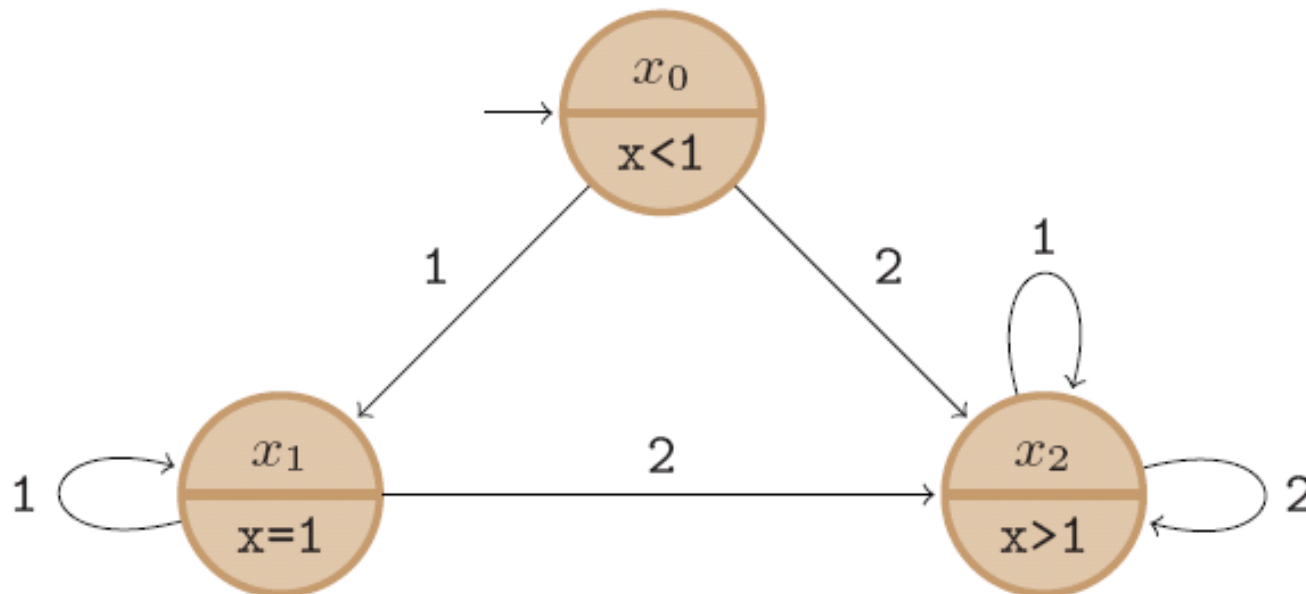
We can formally model software as a **LTS**. The states are all the possible memory configurations and the transition relation describes how the memory contents are changed by the execution of instructions.

Suppose that we want to compute the average of a stream of numbers but we do not know a priori the length of the stream. One possible way to compute the average is to update the average upon the reception of a new number in the stream.

```
 $x := 0;$   
 $n := 0;$   
while true do  
     $y := \text{read}(\text{input});$   
     $x := x \frac{n}{n+1} + y \frac{1}{n+1};$   
     $n := n + 1;$   
end
```

Taken from Tabuada, Verification and Control of Hybrid Systems, Springer Verlag, 2009.

Assume that we are interested in knowing if x is smaller, equal, or greater than 1 when y is restricted to assume values in the set $\{1,2\}$. One possible finite-state model capturing the dynamics of x is represented in the following:



Taken from Tabuada, Verification and Control of Hybrid Systems, Springer Verlag, 2009.

A unifying framework: continuous processes

Given a control system Σ :

$$\dot{x} = f(x, u) \quad x \in \mathbb{R}^n, u \in \mathbb{R}^m$$

we can define the following LTS

$$T(\Sigma) = (Q, L, \xrightarrow{\quad}, O, H),$$

where:

- $Q = \mathbb{R}^n$
- L is the collection of control signals $u : \mathbb{R} \rightarrow \mathbb{R}^m$
- $p \xrightarrow{u} q$, if $x(\tau, p, u) = q$ for some $\tau \geq 0$
- $O = \mathbb{R}^n$
- H is the identity function

$T(\Sigma)$ captures all information contained in Σ

... by using similar arguments I can associate a LTS to a hybrid system!

A unifying framework: continuous processes

Given a control system Σ :

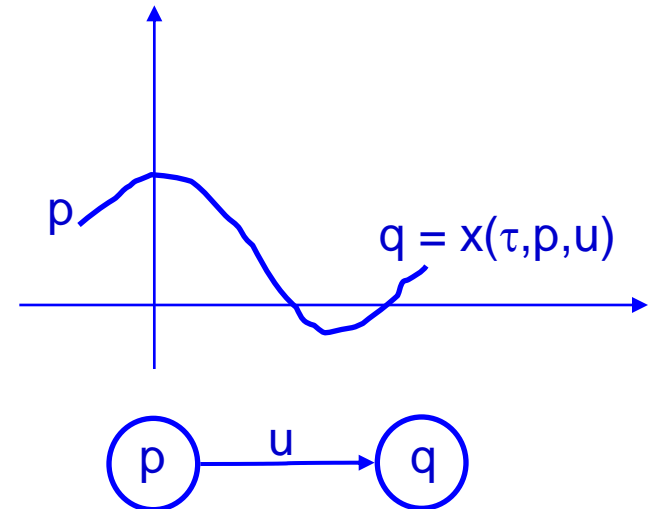
$$\dot{x} = f(x, u) \quad x \in \mathbb{R}^n, u \in \mathbb{R}^m$$

we can define the following LTS

$$T(\Sigma) = (Q, L, \xrightarrow{\quad}, O, H),$$

where:

- $Q = \mathbb{R}^n$
- L is the collection of control signals $u : \mathbb{R} \rightarrow \mathbb{R}^m$
- $p \xrightarrow{u} q$, if $x(\tau, p, u) = q$ for some $\tau \geq 0$
- $O = \mathbb{R}^n$
- H is the identity function



$T(\Sigma)$ captures all information contained in Σ

... by using similar arguments an LTS can be associated to a hybrid system!

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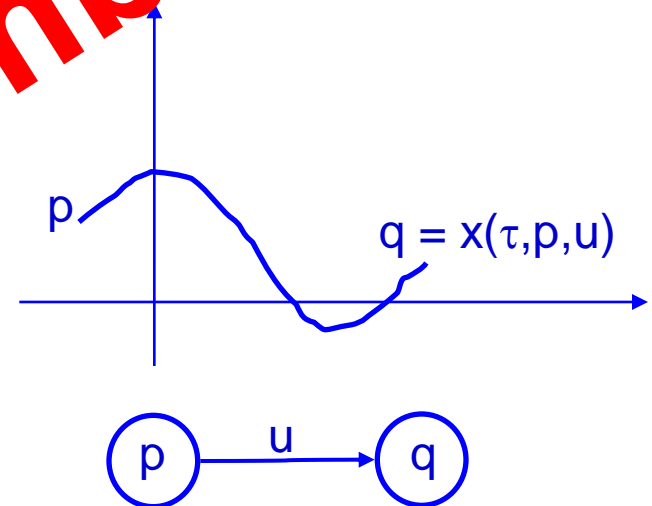
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... by using similar arguments I can associate a LTS to a hybrid system!



$T(\Sigma)$ is not symbolic!

Software modeled by LTS

States are all memory configurations and transition relation describes how memory contents change with execution of instructions

Nonlinear control systems modeled by LTS

Given the nonlinear control system

$$P : dx/dt = f(x,u), x \in \mathbb{R}^n, u \in \mathbb{R}^m$$

consider the LTS

$$T(P) = (Q, L, \xrightarrow{\quad}, O, H),$$

where:

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